
Introduction to Cryptography

Today's homework should be ready soon!
(just got sent to printing)

PCMI 2022 - Undergraduate Summer School

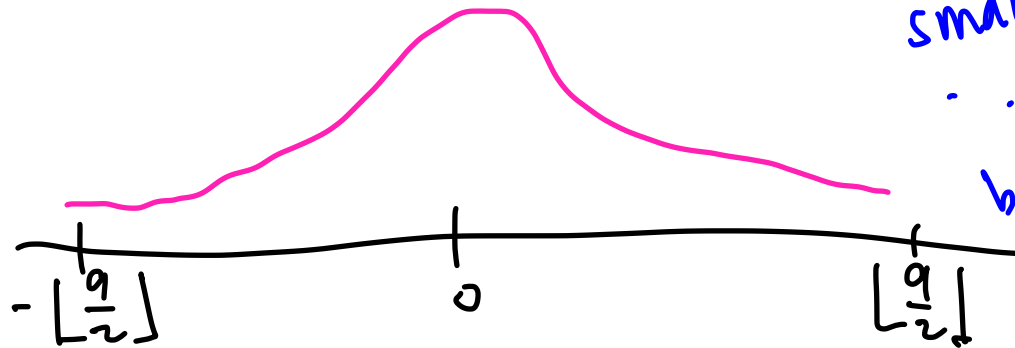
Recall an LWE pair looks like

$$(\vec{a}_i, b_i)$$

• $\vec{a}_i \in \mathbb{F}_q^n$ q prime integer, odd

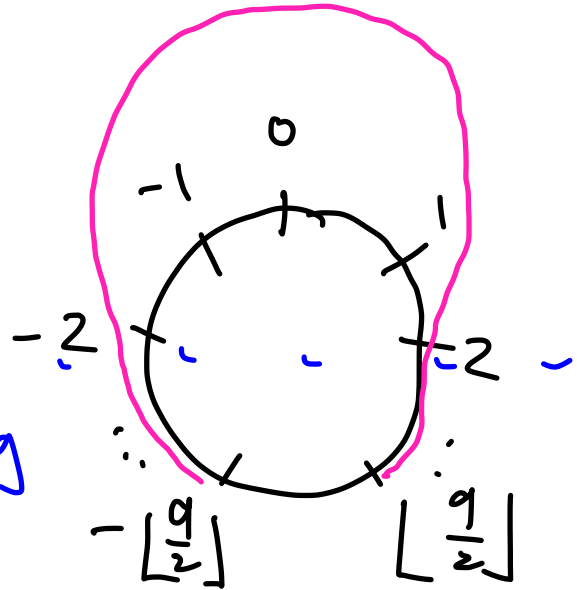
$$b_i = \vec{a}_i \cdot \vec{s} + e_i \in \mathbb{F}_q$$

← "small error"



small

big



Regev's LWE cipher, fix n , prime q , σ^2 variance

key generation

- choose random $\vec{s} \in \mathbb{F}_q^n$, this is the private key
- choose random $\vec{a}_i \in \mathbb{F}_q^n$, errors $e_i \in \mathbb{F}_q$
uniformly chosen from a Gaussian with variance σ^2

• public key are the LWE pairs

$$(\vec{a}_i, b_i = \vec{a}_i \cdot \vec{s} + e_i) \quad i=1, \dots, m$$

Regen in 2005 suggested for fixed n

- $n^2 \leq q_f \leq 2n^2$

- $m = (1 + \varepsilon)(n+1) \log q_f$ for any $\varepsilon > 0$

- $\sigma = \frac{q_f}{\sqrt{2\pi n} (\log n)^2}$

Regen LWE encryption

• Choose a random subset $T \subseteq \{1, \dots, m\}$

• Compute $\vec{a} = \sum_{i \in T} \vec{a}_i$

• To send $x=0$, $b = \sum_{i \in T} b_i$

To send $x=1$, $b = \sum_{i \in T} b_i + \lfloor \frac{q}{2} \rfloor$

• Send (\vec{a}, b)

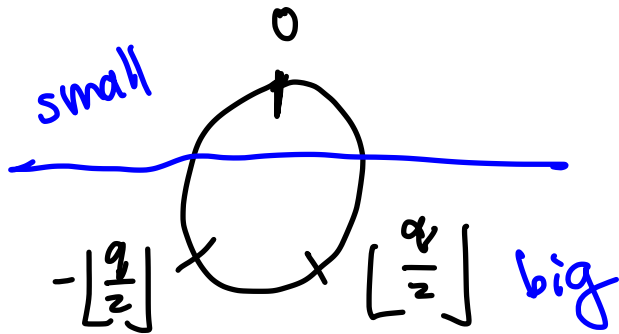
← biggest number

Regen LWE decryption

Compute $\vec{a} \cdot \vec{s} - b = \begin{cases} \sum_{i \in T} e_i & \text{if } x=0 \\ \sum_{i \in T} e_i + \lfloor \frac{q}{2} \rfloor & \text{if } x=1 \end{cases}$

With high probability, $x=0$ if $\vec{a} \cdot \vec{s} - b$ is "small"

$x=1$ if $\vec{a} \cdot \vec{s} - b$ is "big"



Next week

· Monday: Fully homomorphic encryption (FHE)

· Tuesday / Thursday: Ring LWE

· Friday: Cryptography in the real world

Algebraic number theory background for RLWE

- A number field K is a field containing \mathbb{Q} and such that $\dim_{\mathbb{Q}} K = n < \infty$

K satisfies properties of a vector space / \mathbb{Q}

- The number n is called the degree of K

- $\alpha \in K$

consider $\{1, \alpha, \alpha^2, \alpha^3, \dots, \alpha^n\}$

degree of K over \mathbb{Q}

This is $n+1$ elements in a vector space of dim n ,
so there must be a relation

$$a_n \alpha^n + a_{n-1} \alpha^{n-1} + \dots + a_1 \alpha + a_0 = 0$$

for $a_i \in \mathbb{Q}$

$$f(x) = a_n x^n + \dots + a_0 \quad \text{then} \quad f(\alpha) = 0$$

• From the fact that α satisfies a polynomial of degree $\leq n$ with coeffs in \mathbb{Q} ,

we can get that there is a unique

monic irreducible polynomial $p \in \mathbb{Q}[x]$

such that $p(\alpha) = 0$.

only trivial factorization over \mathbb{Q}

$a_n = 1$, leading coefficient

This unique poly is the minimal polynomial of α .

- Inside of K , take the set of elements α such that the minimal polynomial of α has coefficients in \mathbb{Z} . This set is a ring, we call it the ring of integers of K .

Ex: $\frac{1+\sqrt{5}}{2}$ is an "integer"

$\frac{1}{2} \in K$ root of $2x-1$ or $x-\frac{1}{2}$

So $\mathbb{Z} \subseteq \mathcal{O}_K$

← ring of integers
"mathcal O sub K"

Example: $K = \mathbb{Q}(i) = \{a+bi, a, b \in \mathbb{Q}, i^2 = -1\}$

$\mathcal{O}_K = \mathbb{Z}[i] = \{a+bi, a, b \in \mathbb{Z}, i^2 = -1\}$

Primitive element theorem (adapted)

If K is a number field of degree n , then there is $\gamma \in K$ such that

$$K = \mathbb{Q}(\gamma) = \left\{ a_0 + a_1\gamma + a_2\gamma^2 + \dots + a_{n-1}\gamma^{n-1} : a_i \in \mathbb{Q} \right\}$$

We call γ a primitive element, and the minimal polynomial of γ has degree n in this case.

Suppose that $K = \mathbb{Q}(\gamma)$ is a number field of degree n , then there are n injective ring homomorphisms

$$K \hookrightarrow \mathbb{C}$$

\leftarrow complex numbers

$$\gamma \mapsto \gamma_1$$

$$\gamma \mapsto \gamma_2$$

\vdots

$$\gamma \mapsto \gamma_n$$

$\left. \begin{array}{l} \gamma_1, \dots, \gamma_n \text{ are the} \\ n \text{ roots of the} \\ \text{minimal polynomial} \\ \text{of } \gamma \end{array} \right\}$

Example $K = \mathbb{Q}(\alpha)$ with $\alpha^2 - 2 = 0$

$$K \hookrightarrow \mathbb{R} \subseteq \mathbb{C}$$

$$\alpha \mapsto \sqrt{2}$$

$$\alpha \mapsto -\sqrt{2}$$

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That's all for now!