# Introduction to Cryptography <br> 99 problems and LWE is one 

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We say that Problem $A$ reduces to Problem $B$ if, given a solution to Problem B, we can Solve Problem A.

Search LWE (Learning With ERrors) Problem Given a prime $q$ and a positive
hold on
Definition: LWE pairs
Given a prime of and a positive integer $n$ form pairs $\left(\vec{a}_{i}, b_{i}\right)$ with $\vec{a}_{i} \in \mathbb{F}_{q}^{n}, b_{i} \in \mathbb{F}_{q}$
in the following way

- the vector $\vec{a}_{i}$ is chosen uniformly at Random from $\mathbb{F}_{q}^{n}$
- $b_{i}=\vec{a}_{i} \cdot \vec{s}+e_{i}$ for $\vec{s}$ a fixed element of $\mathbb{F}_{q}^{n}$ and $e_{i}$ a "small" Random element of $\mathbb{F}_{q}$.
Pairs like this $\uparrow$ are $L W E_{q, \vec{s}, x}$ pairs

What is X, or what is "small":
We usually use $X$ which is a discrete Gaussian/normal distribution constrained by $-\frac{9}{2}<x<\frac{9}{2}$


LWE pairs: secret $\vec{s}$

$$
\left(\overrightarrow{a_{i}}, \quad b_{i}=\overrightarrow{a_{i}} \cdot \vec{s}+e_{i}\right)
$$

Search LWE Problem
Given a certain number of LWE pairs $\left(\vec{a}_{i}, b_{i}\right)$, find $\vec{s}$.

Decision LWE problem
Given some number of pairs $\left(\overrightarrow{a_{i}}, b_{i}\right)$ determine if they are LWE or if the $b_{i} s$ were chosen at Random (separately from the $\vec{a}_{i} s$ )

$$
4 \quad 1 \quad 2
$$



$$
\begin{aligned}
& \operatorname{PR}(X=0)=\frac{4}{7} \\
& \operatorname{PR}(X=1)=\frac{1}{7} \\
& \operatorname{PR}(X=2)=\frac{2}{7}
\end{aligned}
$$

Theorem:
(1) The decision LWE problem reduces (in polynomial time) to the search LWE problem.
(2) If $q$ is polynomial in $n$, the search LWE problem reduces to the decision LWE problem.
(1) Given pairs $\left(\overrightarrow{a_{i}}, b_{i}\right)$
put them in the search LWE solver to get $\vec{s}$
Then check if $\quad \vec{a}_{i} \cdot \vec{s}-b_{i}$ is distributed like a Gaussian
(2) Given LWE pairs $\left(\vec{a}_{i}, b_{i}\right)$

We can ${ }^{\text {check }}$ guess the first coordinate of $\vec{s}$
in the following way:
Suppose we guess it's $g \in \mathbb{F}_{9}$
For each $i$, choose $r_{i} \in \mathbb{F}_{q}$ at random and form the pair

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(\underbrace{\overrightarrow{a_{i}}+\left(r_{i}, 0,0 \ldots 0\right)}_{\text {new } \vec{a}_{i}^{\prime}}, b_{i}^{\prime}+g r_{i})
$$

Feed the new pairs in the decision LWE solver If the pairs are LWE then the guess is correct. If they are not, guess again.
@home, check why.

In 2005, Reger gave a quantum reduction of the "GapSUP" to the search LWE problem.

Later on, Peikent gave a Classical Reduction of the GapSUP problem for large 9 $\left(q \geqslant 2^{n / 2}\right)$ to the search LWE problem.

In 2008, Regex showed that if $q$ is a product of small primes + erRor is Gaussian, then GapSUP reduces to search
LW.

Definition Short integer solution $S I S_{\beta}$
Fix $\beta>0$, of prime. Given an $n \times m$ matrix
$A$ with entries in $\mathbb{F}_{q}$ find $\vec{z} \neq 0, \vec{z} \in \mathbb{Z}^{m}$
such that

$$
\begin{aligned}
& \|\vec{z}\| \leq \beta \\
& A \vec{z} \equiv 0 \bmod q
\end{aligned}
$$

decision LWE Reduces to SIS.
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## That's all for now!

