THE UNIVERSITY OF VERMONT COLLEGE OF ENGINEERING \& MATHEMATICAL SCIENCES

## QUICKSORT



CS 124 / Department of Computer Science

## Quicksort

Like merge sort, quicksort is a "divide-and-conquer" algorithm.
Divide-and-conquer algorithms break a problem down recursively and then combine the results of subproblems to produce the final result.

## Recursion

- Recursive functions call themselves.
- Recursive functions require a base case or base cases to prevent infinite digress.
- Examples seen earlier in the course:
- Merge sort
- Fibonacci
- Factorial
- Shamos' subsequence sum


## Quicksort

// pseudo-code (see: Stephens, p. 145)
quicksort(vector, start, end) pick an element to divide the array (a.k.a. "pivot") move items less than the pivot to the left of the pivot move items greater than or equal to the pivot to the right of the pivot let $p$ be the index of the pivot quicksort(vector, start, p-1) quicksort(vector, $p+1$, end)

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## Quicksort

## Animation



Animation source: Wikipedia

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This is the basic idea, but there are lots of implementation details to consider!

## Quicksort

partition(vector, start, end)
pivot $=$ vector[end]
i = start
for j from start to end
if vector[j] < pivot then swap vector[i] and vector[j]
$i=i+1$
swap vector[i] and vector[end] return i
quicksort(vector, start, end)
if start < end then
$p=$ partition(vector, start, end)
quicksort(vector, start, p-1)
quicksort(vector, $p+1$, end)

## Quicksort

```
partition(vector, start, end)
    pivot = vector[end]
    i = start
    for j from start to end
        if vector[j] < pivot then
        swap vector[i] and vector[j]
        i = i + 1
    swap vector[i] and vector[end]
    return i
quicksort(vector, start, end)
    if start < end then
        p = partition(vector, start, end)
        quicksort(vector, start, p - 1)
        quicksort(vector, p + 1, end)
```


## Quicksort

partition(vector, start, end)
pivot $=$ vector[end]
i = start
for j from start to end
if vector[j] < pivot then
swap vector[i] and vector[j]
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swap vector[i] and vector[end]

$$
\text { pivot }=9
$$

return i
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\text { pivot }=5
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return i
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$$
\text { pivot }=4
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return i
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## Quicksort

Quicksort has many different implementations and performance varies significantly with implementation.

- How do we choose a pivot (divider)? There are many schemes.
- Quicksort doesn't do well on sorted lists or when there are many repeated values. Quicksort can have worst-case complexity of $O\left(\mathrm{n}^{2}\right)$.
- In the best cases, with a good implementation, quicksort outperforms merge sort (and heap sort, which we'll see a little later).
- Sometimes it is implemented in a hybrid form, with insertion sort used when subproblems become small.


## Merge sort

Recursion depth = log (n)


## Quicksort

## Typical case

$$
\text { Recursion depth } \cong \log (n)
$$



$$
\begin{aligned}
& \square \square \square \square \square \square \square \square \square \square \square \\
& \square \square \square \square \square \square \square \square \square \square \square \square \\
& \square \square \square \square \square \square \square \square \square \square \square \square \square \square \square \square
\end{aligned}
$$

## Quicksort

## Worst case

$$
\begin{aligned}
& \begin{array}{|l|l|l|l|l|l|l|l|l|l|l|l|l|l|l|l|}
\hline \text { a } & \text { b } & \text { c } & \text { d } & \text { e } & \text { f } & \text { g } & \text { h } & \text { i } & \text { j } & \text { k } & \text { l } & \text { m } & \text { n } & \text { o } & \text { p } \\
\hline
\end{array} \\
& \begin{array}{l|l|l|l|l|l|l|l|l|l|l|l|l|l|l|}
\hline \mathrm{a} & \mathrm{~b} & \mathrm{c} & \mathrm{~d} & \mathrm{e} & \mathrm{f} & \mathrm{~g} & \mathrm{n} & \mathrm{i} & \mathrm{j} & \mathrm{k} & \mathrm{I} & \mathrm{~m} & \mathrm{n} & \mathrm{o} \\
\hline
\end{array} \\
& \begin{array}{l|l|l|l|l|l|l|l|l|l|l|l|l|l|}
\hline a & \mathrm{~b} \\
\mathrm{c} & \mathrm{c} & \mathrm{~d} & \mathrm{e} & \mathrm{f} & \mathrm{~g} & \mathrm{~h} & \mathrm{i} & \mathrm{j} & \mathrm{k} & \mathrm{I} & \mathrm{~m} & \mathrm{n} & \mathrm{o} \\
\hline
\end{array}
\end{aligned}
$$

## Quicksort

## Choosing a pivot

| Method | Pro | Con |
| :--- | :--- | :--- |
| Choose element at one <br> end of the vector | Works pretty well in most cases; super <br> easy implementation | Can be horrible with sorted or <br> partially sorted list |
| Pick at random | Works pretty well in most cases | Cost of choosing at random; not <br> really worth the trouble |
| Median-of-three | Works pretty well in most cases and <br> not too costly; reduces likelihood of <br> worst-case performance | A little overhead (three extra <br> comparisons and up to three extra <br> swaps) |

## Quicksort

## Median-of-three

// pseudo-code, given some vector v , and start and end indices
middle $=($ start + end $) / 2 ;$
if element indexed by middle is less than element indexed by start swap middle and start
if element indexed by end is less than element indexed by start swap end and start
if element indexed by middle is less then element indexed by end swap middle and end
let the pivot be the element now at the end position

## Quicksort

## Median-of-three

## Quicksort

## Median-of-three



## Quicksort

Median-of-three

## Quicksort

Median-of-three


## Quicksort

Median-of-three
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## Quicksort

Median-of-three

## Quicksort

Median-of-three

## Quicksort

Median-of-three

$$
\begin{aligned}
& \text { pivot }
\end{aligned}
$$

## Quicksort Complexity

The partition function must process $O(n)$ elements, and the recursion depth is typically $O(\log n)$, hence $O(n \log n)$. This is the best case and average case.

The worst case is, as we have seen, when we have a linear chain of partitions, and a recursion depth of $O(n)$. Again, the partition function processes $O(n)$ elements so we have $O\left(n^{2}\right)$.

However, we can avoid worst-case performance in many situations by choosing our pivot well. This is why many implementations of quicksort use the median-of-three approach, or something similar.

## Quicksort Stability

Is quicksort stable? MAYBE

## Quicksort hybrid

It is not uncommon that quicksort is hybridized with another algorithm -typically insertion sort.

When we have small vectors, we've seen that swapping an element with itself is not uncommon. One way to address this is to use insertion sort for vectors below a certain size, say around ten elements. This gives quicksort a modest speedup in many cases.

## Comparison

| Algorithm | Time <br> complexity | Space <br> complexity | Stable | Comment |
| :---: | :---: | :---: | :---: | :---: |
| Bubble sort | $O\left(n^{2}\right)$ | $O(1)$ | yes | can tell if list is already sorted |
| Selection <br> sort | $O\left(n^{2}\right)$ | $O(1)$ | no | performs fewest swaps |
| Insertion sort | $O\left(n^{2}\right)$ | $O(1)$ | yes | ignores unsorted portion of vector <br> $/$ can process data on-line |
| Merge sort | $O(n$ log $n)$ | $O(n)$ | yes | recursive divide-and-conquer |
| Quicksort | $O(n \log n)$ | $O(1)$ | no | recursive divide-and-conquer |

## Comparison



## Summary

- Quicksort is a recursive divide-and-conquer algorithm.
- Quicksort has $O(n \log n)$ time complexity.
- Quicksort has $O(1)$ space complexity.
- Quicksort in its original design is not stable, but at the expense of increasing space complexity it can be implemented as a stable algorithm.

